

On Simulating Finite Markov Chains by the Splitting and Roulette Approach

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The splitting and Russian roulette approach, suggested by von Neumann [1], is one of possible ways to increase the efficiency of stochastic computer simulation of random processes. The idea of the approach consists in a parallel simulation of several independent paths of a random process corresponding to the behavior of a complex system number investigation. To each path procedures of splitting and roulette are applied. In most versions described in the literature (see, f.e., [2] or [3]) strong restrictions on transition probabilities are assumed. However, in [4] a more general technique was introduced and a corresponding mathematical theory was developed. The present paper is devoted to a brief review of the theory for the case of simulating finite Markov chains. We also suggest an illustrating example.

Let us consider a finite Markov chain with the state space $\Omega = \{0, 1, \dots, m\}$, a transition matrix P and a stationary distribution $\pi = (\pi_0, \dots, \pi_m)^T$. Assume that for an integer number N all elements of the matrix P^N are positive. The problem is to evaluate the vector π by simulation.

As the immediate simulation method we will consider the well known regenerative approach [5]. Let $\{i^*\}$ be the initial state. The method consists in the simulation of several paths of the chain beginning in $\{i^*\}$ till return to this state. As the estimator of π_i ($i = 0, 1, \dots, m$) the number of hits to the state $\{i\}$ by all paths divided by the number of all steps by all paths is used.

To describe the branching technique which is a formalization of the splitting and roulette approach let us introduce the notion of experimental design.

Let β_0, \dots, β_m be arbitrary fixed nonnegative numbers and it is possible that $\beta_i = 0$ if the state $\{i\}$ is a state in which every path will be cancelled.

When a path transits from a state i to a state j , $i \in \Omega$, $j \in \Omega \setminus \{i^*\}$ when $\beta_j \leq \beta_i$ we will simulate additionally η paths beginning in the state j where

$$\eta = \begin{cases} \lfloor \beta_j / \beta_i \rfloor - 1 & \text{with a probability } 1 - \alpha \\ \lfloor \beta_j / \beta_i \rfloor & \text{with a probability } \alpha, \end{cases}$$

$\alpha = \beta_j / \beta_i - \lfloor \beta_j / \beta_i \rfloor$ and $\lfloor a \rfloor$ designates the integer part of a . When $\beta_j < \beta_i$ we will cancel the current path with the probability β_j / β_i . As the estimator of π_k ($k = 0, 1, \dots, m$) we will take the number of hits into the state k by all path multiplied by β_k and divided by the sum of all such numbers. One of β_k could be equal to 0 since $\pi_k = 1 - \sum_{i \neq k} \pi_i$ and one of the π_i could be recalculated by others.

We will call experimental design the discrete probability measure $\tau = \{\tau_0, \dots, \tau_m\}$, $\tau_i \geq 0$ ($i = 0, \dots, m$), $\sum \tau_i = 1$, where

$$\tau_i = \beta_i \pi_i / \sum_{k=0}^m \beta_k \pi_k.$$

In [4] it was proved that the estimators are asymptotically unbiased and the covariance matrix of the estimators of (π_1, \dots, π_m) multiplied by all steps is approximately equal

$$D(\tau) = W^T B(\tau) W, \quad W = (I - P_-)^{-1},$$

where P_- is the matrix P with the first row and the first column rejected,

$$B(\tau) = \sum_{k=0}^{m-1} \frac{\pi_k^2}{\tau_k} B_k, \quad B_k = (p_{ki}\delta_{ij} - p_{ki}p_{kj})_{i,j=1}^m,$$

$$\delta_{ij} = 1 \text{ for } i = j \text{ and } \delta_{ij} = 0 \text{ for } i \neq j.$$

The immediate simulation method is a particular case when $\beta_0 = \beta_1 = \dots = \beta_m$ that is $\tau = \pi$. The natural efficiency criterion for the problem is the determinant of $D(\tau)$. Since W does not depend on τ the criterion is reduced to $\det B(\tau)$.

A design τ^* will be called D -optimal design if it minimizes the magnitude of $\det B(\tau)$.

Consider a particular case of the Markov chain embedded into the random process, corresponding to the length of queue with one server and m places for waiting. We will consider the simplest case when the input stream is a simplest stream and the time of service is an exponentially distributed random value.

Let ρ be the load of the system. Then the matrix P is of the form

$$P = \begin{pmatrix} 0 & 1 & 0 & 0 & \dots \\ 1 - \Delta & 0 & \Delta & 0 & \\ 0 & 1 - \Delta & 0 & \Delta & \\ \vdots & & & & \\ & & & 1 - \Delta & 0 & \Delta \\ & & & 0 & 1 - \Delta & \Delta \end{pmatrix}, \Delta = \rho / (1 + \rho),$$

$$\pi_1 = \Delta\pi_0, \pi_{i+1} = \rho^i \Delta \pi_0, i = 1, \dots, m - 1, \pi_0 = 1 / (1 + \Delta + \Delta\rho + \dots + \Delta\rho^{m-1}).$$

We can calculate by induction that

$$\det B(\tau) = \left(\prod_{i=1}^m \frac{\pi_i^2}{\tau_i} \right) \Delta^m (1 - \Delta)^m.$$

From here we found that the D -optimal design is $\tau^* = \{0, \frac{1}{m}, \dots, \frac{1}{m}\}$ and

$$I(\tau) = \left(\frac{\det B(\pi)}{\det B(\tau)} \right)^{1/m} = \left(\prod_{i=1}^m \pi_i \right)^{-1/m} / m$$

with $\tau = \tau^*$. Note that $I(\tau)$ is a natural measure of efficiency of the design τ . It means the ratio of the number of steps needed by immediate simulation for obtaining results with a given accuracy to the respective number for the splitting and roulette approach. In table 1 values of the efficiency are given.

Table 1. Values of I .

m	5	5	5	10	10	10
ρ	1/2	1/3	1/4	1/2	1/3	1/4
I	2.0	4.0	6.8	6.0	31.6	109

We can conclude that the technique is very powerful in this case. For practical implementation the optimal values $\beta_j = 1/\pi_j, j \neq 0$ could be evaluated through the current simulation results in the style of sequential approach. As the initial state we can take $\{i^*\} = \{1\}$.

We can offer to use the design τ^* for the more complicated chains, simulation whereas the efficiency could be evaluated numerically.

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